## PDEs and Graph Based Learning

Summer School on Random Structures in Optimizations and Related Applications

## Lecture 1: $k$-means and spectral clustering Instructor: Jeff Calder (jcalder@umn.edu)

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Clustering


## $k$-means clustering

Let $x_{1}, x_{2}, \ldots, x_{m}$ be datapoints in $\mathbb{R}^{n}$. The $k$-means algorithm is guided by the task of minimizing the energy over the choice of cluster centers $c_{i}$

$$
E\left(c_{1}, c_{2}, \ldots, c_{k}\right)=\sum_{i=1}^{m} \min _{1 \leq j \leq k}\left\|x_{i}-c_{j}\right\|^{2}
$$

Minimizing $E$ is an NP-hard problem.


## $k$-means clustering

$k$-means algorithm: We start with some randomized initial values for the means $c_{1}^{0}, c_{2}^{0}, \ldots, c_{k}^{0}$, and iterate the steps below until convergence.

1. Update the clusters

$$
\begin{equation*}
\Omega_{j}^{t}=\left\{x_{i}:\left\|x_{i}-c_{j}^{t}\right\|^{2}=\min _{1 \leq \ell \leq k}\left\|x_{i}-c_{\ell}^{t}\right\|^{2}\right\} \tag{1}
\end{equation*}
$$

2. Update the cluster centers (means)

$$
\begin{equation*}
c_{j}^{t+1}=\frac{1}{\# \Omega_{j}^{t}} \sum_{x \in \Omega_{j}^{t}} x . \tag{2}
\end{equation*}
$$

The algorithm converges when $c_{j}^{t+1}=c_{j}^{t}$ for all $j$.
$k$-means clustering result


## Poor clustering by $k$-means

Clustering depends on the random initialization.

$k$-means clustering in Python (.ipynb)

Lemma on centroids
Lemma 1. Let $y_{1}, y_{2}, \ldots, y_{m}$ be points in $\mathbb{R}^{n}$, and define the function $f: \mathbb{R}^{n} \rightarrow \mathbb{R}$ by

$$
f(x)=\sum_{i=1}^{m}\left\|y_{i}-x\right\|^{2}
$$

Then the unique minimizer of $f$ is the centroid

$$
c=\frac{1}{m} \sum_{i=1}^{m} y_{i}
$$

In particular, $f(c)<f(x)$ if $x \neq c$.

$$
\begin{aligned}
\text { Proof }: f(c) & =\sum_{i=1}^{m}\left\|y_{i}-c\right\|^{2} \\
& =\sum_{i=1}^{m}\left(\left\|y_{i}\right\|^{2}-2 c^{\top} y_{i}+\|c\|^{2}\right)
\end{aligned}
$$

$$
\begin{aligned}
& =\sum_{i=1}^{m}\left\|y_{i}\right\|^{2}-2 c^{T} \underbrace{m}_{=m c} y_{i=1}^{m}+\sum_{i=1}^{m}\|c\|^{2} \\
& =\sum_{i=1}^{m}\left\|y_{i}\right\|^{2}-2 c^{\top} m c+m\|c\|^{2} \\
& =\sum_{i=1}^{m}\left\|y_{i}\right\|^{2}-m\|c\|^{2} \\
& f(c)=\sum_{i=1}^{m}\left(\left\|y_{i}\right\|^{2}-\|c\|^{2}\right) \\
& f(x)=\sum_{i=1}^{m}\left\|y_{i}-x\right\|^{2} \\
& =\sum_{i=1}^{m}\left(\left\|y_{i}\right\|^{2}-2 x y_{y_{i}}+\|x\|^{2}\right)
\end{aligned}
$$

$$
\begin{aligned}
& =\sum_{i=1}^{m}\left(\left\|y_{i}\right\|^{2}-\|c\|^{2}\right)+\sum_{i=1}^{m}\left(\|c\|^{2}-2 x^{\top} y_{i}+\|x\|^{2}\right) \\
& =f(c)+m\|c\|^{2}-2 x^{\top} m c+m\|x\|^{2} \\
& =f(c)+m\left(\|c\|^{2}-2 x^{\top} c+\|x\|^{2}\right) \\
& =f(c)+m\|c-x\|^{2} \\
f(x) & =f(c)+m\|c-x\|^{2}
\end{aligned}
$$

A1t. Proff $\nabla f(x)=2 \sum_{i=1}^{m}\left(y_{i}-x\right)=0$

$$
\ldots y_{i}=c
$$

IIIII

## Convergence of $k$-means

Recall the $k$-means enegy

$$
\begin{equation*}
E\left(c_{1}, c_{2}, \ldots, c_{k}\right)=\sum_{i=1}^{m} \min _{1 \leq j \leq k}\left\|x_{i}-c_{j}\right\|^{2} \tag{3}
\end{equation*}
$$

Theorem 2. The $k$-means algorithm descends on the energy (3), that is

$$
\begin{equation*}
E\left(c_{1}^{t+1}, c_{2}^{t+1}, \ldots, c_{k}^{t+1}\right) \leq E\left(c_{1}^{t}, c_{2}^{t}, \ldots, c_{k}^{t}\right) \tag{4}
\end{equation*}
$$

Furthermore, we have equality in (4) if and only if $c_{j}^{t+1}=c_{j}^{t}$ for $j=1, \ldots, k$, and hence the $k$-means algorithm converges in a finite number of iterations.

## Note:

- $k$-means does not in general find a global minimum of $E$.
- It is useful because it is fast, guaranteed to converge, and often finds good clustering.


## $k$-means on two-moons



- Sometimes a single point is not a good representative of a cluster, in Euclidean distance.
- Instead, we can try to cluster points so that nearby points are assigned to the same cluster, without specifying cluster centers.


## Weight matrix

Let $x_{1}, x_{2}, \ldots, x_{m}$ be points in $\mathbb{R}^{n}$. To encode which points are nearby, we construct a weight matrix $W$, which is an $m \times m$ symmetric matrix where $W(i, j)$ represents the similarity between datapoints $x_{i}$ and $x_{j}$. A common choice for the weight matrix is Gaussian weights

$$
\begin{equation*}
W(i, j)=\exp \left(-\frac{\left\|x_{i}-x_{j}\right\|^{2}}{2 \sigma^{2}}\right) \tag{5}
\end{equation*}
$$

where the $\sigma$ is a free parameter that controls the scale at which points are connected.

## Graph cuts for binary clustering

A graph-cut approach to clustering minimizes the graph cut energy

$$
\begin{equation*}
E(z)=\frac{1}{2} \sum_{i=1}^{m} \sum_{j=1}^{m} W(i, j)|z(i)-z(j)|^{2} \tag{6}
\end{equation*}
$$

over label vectors $z \in\{0,1\}^{m}$.

Notes:

- The value $z(i) \in\{0,1\}$ indicates which cluster $x_{i}$ belongs to.
- The graph-cut energy is the sum of the edge weights $W(i, j)$ that must be cut to separate the dataset into two clusters.


## Balanced graph cuts for binary clustering

Minimizing the graph cut energy

$$
E(z)=\frac{1}{2} \sum_{i=1}^{m} \sum_{j=1}^{m} W(i, j)|z(i)-z(j)|^{2}
$$

can lead to very unbalanced clusters (e.g., one cluster can have just a single point).

A more useful approach is to minimize a balanced graph cut energy

$$
\begin{equation*}
E_{\text {balanced }}(z)=\frac{\frac{1}{2} \sum_{i=1}^{m} \sum_{j=1}^{m} W(i, j)|z(i)-z(j)|^{2}}{\sum_{i=1}^{n} z(i) \sum_{j=1}^{n}(1-z(j))} \tag{7}
\end{equation*}
$$

The denominator is the product of the number of points in each cluster, which is maximized when the clusters are balanced.

Balanced graph-cut problems are NP hard.

## Relaxing the graph cut problem

To relax the graph-cut problem, we consider minimizing the graph cut energy

$$
E(z)=\frac{1}{2} \sum_{i=1}^{m} \sum_{j=1}^{m} W(i, j)|z(i)-z(j)|^{2}
$$

over all real-vectors $z \in \mathbb{R}^{m}$. We still have a balancing issue (here $z=0$ is a minimizer), so we impose the balancing constraints

$$
\text { (\$) } \mathbf{1}^{T} z=\sum_{i=1}^{m}\left(z_{i}=0 \quad \text { and } \quad\|z\|^{2}=\sum_{i=1}^{m} z(i)^{2}=1\right.
$$

Definition 3. The binary spectral clustering problem is
Minimize $E(z)$ over $z \in \mathbb{R}^{m}$, subject to $\mathbf{1}^{T} z=0$ and $\|z\|^{2}=1$.
The resulting clusters are $C_{1}=\left\{x_{i}: z(i)>0\right\}$ and $C_{2}=\left\{x_{i}: z(i) \leq 0\right\}$.

## The graph Laplacian and Fiedler vector

Let $W$ be a symmetric $m \times m$ matrix with nonnegative entries.
Definition 4. The graph Laplacian matrix $L$ is the $m \times m$ matrix

$$
\begin{equation*}
L=D-W \tag{8}
\end{equation*}
$$

where $D$ is the diagonal matrix with diagonal entries

$$
D(i, i)=\sum_{j=1}^{m} W(i, j) .=\text { degree ef } i
$$

Lemma 5. Then the graph cut energy can be expressed as

$$
E(z)=\frac{1}{2} \sum_{i=1}^{m} \sum_{j=1}^{m} W(i, j)|z(i)-z(j)|^{2}=z^{T} L z
$$

where $L$ is the graph Laplacian.

Proof: $E(z)=\frac{1}{2} \sum_{i=1}^{m} \sum_{j=1}^{m} w(i j)\left(z(i)^{2}-2 z(i) z(j)+z(j)^{2}\right)$

$$
\begin{aligned}
= & \frac{1}{2} \sum_{i=1}^{m} \sum_{j=1}^{m} w(i, j) z(i)^{2}+\frac{1}{2} \sum_{i=1}^{m} \sum_{j=1}^{m} w(i j) z(j)^{2} \\
& -\sum_{i=1}^{m} \sum_{j=1}^{m} w(i, j) z(i) z(j) \\
= & \sum_{i=1}^{m} \underbrace{(W z)(i)}_{\sum_{j=1}^{m} w(i, j) z(i)^{2}-\sum_{i=1}^{m} \sum_{j=1}^{m} w(i, j) z(j) z(i)} \\
= & \sum_{i=1}^{m} D(i, i) z(i)^{2}-\sum_{i=1}^{m}(w z)(i) z(i)
\end{aligned}
$$

$$
=z^{\top} D z-z^{\top} \omega z=z^{\top} L z
$$

since $L=D-\omega$. $\square$

Properties of the graph Laplacian
Lemma 6. Let $L=D-W$ be the graph Laplacian corresponding to a symmetric matrix $W$ with nonnegative entries. The following properties hold.
(i) $L$ is symmetric.
(ii) $L$ is positive semi-definite (i.e., $z^{T} L z \geq 0$ for all $z \in \mathbb{R}^{m}$ ).
(iii) All eigenvalues of $L$ are nonnegative, and the constant vector $z=\mathbf{1}$ is an eigenvector of $L$ with eigenvalue $\lambda=0$.

$$
\begin{aligned}
& L 1=D 1-w 1 \\
& (w 1)(i)=\sum_{j=1}^{m} w(i, j) 1=D(i, i) \\
& \Rightarrow W 1=D 1
\end{aligned}
$$

## Fiedler vector

Let $v_{1}, v_{2}, \ldots, v_{m}$ be the eigenvectors of the graph Laplacian, with corresponding eigenvalues

$$
0=\lambda_{1} \leq \lambda_{2} \leq \cdots \leq \lambda_{m} .
$$

Definition 7. The second eigenvector $v_{2}$ of the graph Laplacian $L$ is called the Fiedler vector.

Theorem 8. The Fiedler vector $z=v_{2}$ solves the binary spectral clustering problem

$$
\text { Minimize } E(z) \text { over } z \in \mathbb{R}^{m} \text {, subject to } \mathbf{1}^{T} z=0 \text { and }\|z\|^{2}=1 \text {. }
$$

Proof: Let $z$ be a minimizer

Write $z=a_{1} v_{1}+a_{2} v_{2}+\cdots+a_{m} v_{m}$

$$
\begin{aligned}
& \left(v_{1}=1 / \sqrt{m}, \quad\left\|v_{i}\right\|=1\right) \\
& 0=I^{\top} z=a_{1} I^{\top} v_{1}+a_{2} I I_{v_{2}}^{0}+\cdots+a_{m} Z^{\top} v_{m}^{0} \\
& =a_{1} \frac{m}{\sqrt{m}} \Rightarrow a_{1}=0 \quad \operatorname{sinc} v_{i}^{\top} v_{j}=0 \text {, i才j } \\
& =a_{1} \frac{m}{\sqrt{m}} \Rightarrow a_{1}=0 \\
& 1=\|z\|^{2}=\sum_{i=1}^{m} a_{i}^{2}=a_{2}^{2}+a_{3}^{2}+\cdots \alpha+a_{m}^{2} \\
& E(z)=z^{\top} L z=z^{\top} L \sum_{i=2}^{m} a_{i} v_{i} \\
& =z^{\top} \sum_{i=2}^{m} a_{i} L v_{i}
\end{aligned}
$$

$$
\begin{aligned}
&=z^{T} \sum_{i=2}^{m} a_{i} \lambda_{i} v_{i} \\
&=\sum_{i=2}^{m} a_{i} \lambda_{i} \underbrace{T v_{i}}_{=a_{i}} \\
&=\sum_{i=2}^{m} \lambda_{i} a_{i}^{2} \\
& E(z)=\sum_{i=2}^{m} \lambda_{i} a_{i}^{2} \geq \lambda_{2} \sum_{i=2}^{m} a_{i}^{2}=\lambda_{2} \\
& \quad \lambda_{2} \leqslant \lambda_{3} \leqslant \lambda_{4} \leq \ldots \leq \lambda_{m}
\end{aligned}
$$

Settim $z=v_{2} \quad\left(a_{2}=1, a_{3}=0, a_{4}=0, \ldots\right)$ yields $E(z)=\lambda_{2}$

## Example


(a) Fiedler vector

(b) Spectral Clustering

Figure 1: (a) The Fiedler vector and (b) spectral clustering on the 2-moons dataset.

## Spectral embeddings

Let $v_{1}, v_{2}, v_{3}, \ldots$ be the normalized eigenvectors of $L$, in order of increasing eigenvalues $0=\lambda_{1} \leq \lambda_{2} \leq \cdots$. The spectral embedding corresponding to $L$ is the map $\Phi: I_{m} \rightarrow \mathbb{R}^{k}$ (recall $I_{m}=\{1,2, \ldots, m\}$ are the indices of our datapoints) given by

$$
\begin{equation*}
\Phi(i)=\left(v_{1}(i), v_{2}(i), \ldots, v_{k}(i)\right) \tag{9}
\end{equation*}
$$

Since the first eigenvector $v_{1}$ is the trivial constant eigenvector, it is also common to omit this to obtain the embedding

$$
\Phi(i)=\left(v_{2}(i), v_{3}(i), \ldots, v_{k+1}(i)\right) .
$$

There are other normalizations of the graph Laplacian that are commonly used, such as the symmetric normalization $L=D^{-1 / 2}(D-W) D^{-1 / 2}$, and the spectral embedding for a normalized Laplacian is defined analagously.

## Spectral embeddings

The intuition behind the spectral embedding is encapsulated in the following simple result.

Proposition 9. If $A \subset I_{m}$ is a disconected component of the graph, which means that $W(i, j)=0$ for all $i \in A$ and $j \in I_{m} \backslash A$, then the indicator function of $A$, denoted $u_{A}$, satisfies

$$
L u_{A}=(D-W) u_{A}=0 .
$$

## Spectral embedding of MNIST


(a) Unnormalized

(b) Normalized

Figure 2: Example of spectral embeddings in the plane $k=2$ of the 0,1 , and 2 digits of the MNIST dataset using the unnormalized $L=D-W$ and symmetric normalized $L=D^{-1 / 2}(D-W) D^{-1 / 2}$ graph Laplacians.

## Spectral clustering with more than 2 classes

Suppose we want to cluster the graph into $k$ clusters. Spectral clustering proceeds as follows:

1. Perform a spectral embedding of the graph into $\mathbb{R}^{k}$ (or sometimes $\mathbb{R}^{d}$ where $d \approx k$ ).
2. Run your favorite clustering algorithm in the embedding space $\mathbb{R}^{d}$, such as $k$-means.

## $k$-nearest neighbor graph

The Gaussian weights

$$
W(i, j)=\exp \left(-\frac{\left\|x_{i}-x_{j}\right\|^{2}}{2 \sigma^{2}}\right)
$$

are not always useful in practice, since the matrix $W$ is dense (all entries are nonzero), and the connectivity length $\sigma$ is the same across the whole graph.

It is more common to use a $k$-nearest neighbor graph. Let $d_{k, i}$ denote the Euclidean distance between $x_{i}$ and its $k^{\text {th }}$ nearest Euclidean neighboring point from $x_{1}, \ldots, x_{m}$. A $k$-nearest neighbor graph uses the weights

$$
W(i, j)= \begin{cases}1, & \text { if }\left\|x_{i}-x_{j}\right\| \leq \max \left\{d_{k, i}, d_{k, j}\right\} \\ 0, & \text { otherwise }\end{cases}
$$

The weights need not be binary, and can depend on $\left\|x_{i}-x_{j}\right\|$, similar to the Gaussian weights. The $k$-nearest neighbor graph weight matrix $W$ is very sparse (most entries are zero), so it can be stored and computed with efficiently.

## Exercises

The following proofs were omitted and are left to review as an exercise.

1. Proof of Theorem 2 (Convergence of $k$-means, Theorem 4.2 in lecture notes)
2. Proof of Lemma 6 (Properties of the graph Laplacian, Lemma 4.8 in lecture notes)
3. Proof of Proposition 9 (Spectral embeddings, Proposition 8.4 in lecture notes)

## Spectral clustering in Python (.ipynb)

