PDEs and Graph Based Learning

Summer School on Random Structures in Optimizations and Related Applications

Lecture 1: k-means and spectral clustering Instructor: Jeff Calder (jcalder@umn.edu)

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Clustering





k-means clustering

Let x_1, x_2, \ldots, x_m be datapoints in \mathbb{R}^n . The k-means algorithm is guided by the task of minimizing the energy over the choice of cluster centers c_i

$$E(c_1, c_2, \dots, c_k) = \sum_{i=1}^m \min_{1 \le j \le k} \|x_i - c_j\|^2.$$

Minimizing E is an NP-hard problem.



k-means clustering

k-means algorithm: We start with some randomized initial values for the means $c_1^0, c_2^0, \ldots, c_k^0$, and iterate the steps below until convergence.

1. Update the clusters

(1)
$$\Omega_j^t = \left\{ x_i : \|x_i - c_j^t\|^2 = \min_{1 \le \ell \le k} \|x_i - c_\ell^t\|^2 \right\}.$$

2. Update the cluster centers (means)

(2)
$$c_j^{t+1} = \frac{1}{\#\Omega_j^t} \sum_{x \in \Omega_j^t} x.$$

The algorithm converges when $c_j^{t+1} = c_j^t$ for all j.

Code Demo

k-means clustering result





Poor clustering by k-means

Clustering depends on the random initialization.



k-means clustering in Python (.ipynb)

Lemma on centroids

Lemma 1. Let y_1, y_2, \ldots, y_m be points in \mathbb{R}^n , and define the function $f : \mathbb{R}^n \to \mathbb{R}$ by

$$f(x) = \sum_{i=1}^{m} \|y_i - x\|^2.$$

Then the unique minimizer of f is the centroid

$$c = \frac{1}{m} \sum_{i=1}^{m} y_i.$$

In particular, f(c) < f(x) if $x \neq c$.

$$\begin{array}{ll} P_{roof:} & f(c) = \sum_{i=1}^{m} \|y_i - c\|^2 \\ & = \sum_{i=1}^{m} (\|y_i\|^2 - 2c^T y_i + \|c\|^2) \\ & = \sum_{i=1}^{m} (\|y_i\|^2 - 2c^T y_i + \|c\|^2) \end{array}$$

$$= \sum_{i=1}^{\infty} ||y_{i}||^{2} - 2cT \sum_{i=1}^{\infty} y_{i} + \sum_{i=1}^{\infty} ||e_{i}|^{2}$$

$$= \sum_{i=1}^{\infty} ||y_{i}||^{2} - 2cTmc + m ||c||^{2}$$

$$= \sum_{i=1}^{\infty} ||y_{i}||^{2} - m ||c_{i}|^{2}$$

$$= \sum_{i=1}^{\infty} (||y_{i}||^{2} - m ||c_{i}|^{2})$$

$$f(x) = \sum_{i=1}^{\infty} (||y_{i}||^{2} - 2xTy_{i} + ||x_{i}|^{2})$$

$$= \sum_{i=1}^{\infty} (||y_{i}||^{2} - 2xTy_{i} + ||x_{i}|^{2})$$

$$= \sum_{i=1}^{\infty} (\|v_{i}\|^{2} - \|c_{i}\|^{2}) + \sum_{i=1}^{\infty} (\|v_{i}\|^{2} - a_{x}Tv_{i} + \|v_{i}\|^{2})$$

$$= f(c) + m \||c_{i}\|^{2} - a_{x}Tc + m\|v_{i}\|^{2})$$

$$= f(c) + m (\|c_{i}\|^{2} - a_{x}Tc + \|v_{i}\|^{2})$$

$$= f(c) + m \||c - x||^{2}$$

$$f(x) = f(c) + m \||c - x||^{2}$$

$$Alt. Proof $\nabla f(x) = 2\sum_{i=1}^{\infty} (v_{i} - x) = 0$

$$v_{i} = C$$$$



Convergence of k-means

Recall the k-means energy

(3)
$$E(c_1, c_2, \dots, c_k) = \sum_{i=1}^m \min_{1 \le j \le k} \|x_i - c_j\|^2.$$

Theorem 2. The k-means algorithm descends on the energy (3), that is

(4)
$$E(c_1^{t+1}, c_2^{t+1}, \dots, c_k^{t+1}) \le E(c_1^t, c_2^t, \dots, c_k^t).$$

Furthermore, we have equality in (4) if and only if $c_j^{t+1} = c_j^t$ for j = 1, ..., k, and hence the k-means algorithm converges in a finite number of iterations.

Note:

- k-means does **not** in general find a global minimum of E.
- It is useful because it is fast, guaranteed to converge, and often finds good clustering.

k-means on two-moons



- Sometimes a single point is not a good representative of a cluster, in Euclidean distance.
- Instead, we can try to cluster points so that nearby points are assigned to the same cluster, without specifying cluster centers.

Weight matrix

Let x_1, x_2, \ldots, x_m be points in \mathbb{R}^n . To encode which points are nearby, we construct a weight matrix W, which is an $m \times m$ symmetric matrix where W(i, j) represents the similarity between datapoints x_i and x_j . A common choice for the weight matrix is Gaussian weights

(5)
$$W(i,j) = \exp\left(-\frac{\|x_i - x_j\|^2}{2\sigma^2}\right),$$

where the σ is a free parameter that controls the scale at which points are connected.

Graph cuts for binary clustering

A graph-cut approach to clustering minimizes the graph cut energy

(6)
$$E(z) = \frac{1}{2} \sum_{i=1}^{m} \sum_{j=1}^{m} W(i,j) |z(i) - z(j)|^2$$

over label vectors $z \in \{0, 1\}^m$.

Notes:

- The value $z(i) \in \{0, 1\}$ indicates which cluster x_i belongs to.
- The graph-cut energy is the sum of the edge weights W(i, j) that must be **cut** to separate the dataset into two clusters.

Balanced graph cuts for binary clustering

Minimizing the graph cut energy

$$E(z) = \frac{1}{2} \sum_{i=1}^{m} \sum_{j=1}^{m} W(i,j) |z(i) - z(j)|^2$$

can lead to very unbalanced clusters (e.g., one cluster can have just a single point).

A more useful approach is to minimize a balanced graph cut energy

(7)
$$E_{balanced}(z) = \frac{\frac{1}{2} \sum_{i=1}^{m} \sum_{j=1}^{m} W(i,j) |z(i) - z(j)|^2}{\sum_{i=1}^{n} z(i) \sum_{j=1}^{n} (1 - z(j))}$$

The denominator is the product of the number of points in each cluster, which is maximized when the clusters are balanced.

Balanced graph-cut problems are NP hard.

Relaxing the graph cut problem

To relax the graph-cut problem, we consider minimizing the graph cut energy

$$E(z) = \frac{1}{2} \sum_{i=1}^{m} \sum_{j=1}^{m} W(i,j) |z(i) - z(j)|^2$$

over all real-vectors $z \in \mathbb{R}^m$. We still have a balancing issue (here z = 0 is a minimizer), so we impose the balancing constraints

(1)
$$\mathbf{1}^T z = \sum_{i=1}^m (z_i) = 0$$
 and $||z||^2 = \sum_{i=1}^m z(i)^2 = 1.$

Definition 3. The binary spectral clustering problem is

Minimize E(z) over $z \in \mathbb{R}^m$, subject to $\mathbf{1}^T z = 0$ and $||z||^2 = 1$.

The resulting clusters are $C_1 = \{x_i : z(i) > 0\}$ and $C_2 = \{x_i : z(i) \le 0\}.$



The graph Laplacian and Fiedler vector

Let W be a symmetric $m \times m$ matrix with nonnegative entries.

Definition 4. The graph Laplacian matrix L is the $m \times m$ matrix

$$(8) L = D - W$$

where D is the diagonal matrix with diagonal entries

$$D(i,i) = \sum_{j=1}^{m} W(i,j)$$
. I degree if i

Lemma 5. Then the graph cut energy can be expressed as

$$E(z) = \frac{1}{2} \sum_{i=1}^{m} \sum_{j=1}^{m} W(i,j) |z(i) - z(j)|^2 = z^T L z,$$

where L is the graph Laplacian.

 $\frac{P_{roof}: E(2) = -\frac{1}{2} \sum_{i=1}^{\infty} \sum_{j=1}^{\infty} W(i_{ij}) (7(i_{j})^{2} - 27(i_{j}) + 7(j_{j})^{2})$ $\left(s = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} \left(\frac{1}{1} \int_{-\infty}^{\infty} \frac{1}{1} \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} \frac{1}{1} \int_{-\infty}^{\infty} \int_{-\infty}^{$ $-\sum_{i=1}^{m}\sum_{j=1}^{m}\omega(ij)Z(i)Z(j)$ $= \sum_{i=1}^{\infty} \sum_{j=1}^{\infty} W(i_{ij}) Z(i)^{2} - \sum_{i=1}^{\infty} \sum_{j=1}^{\infty} W(i_{ij}) Z(j) Z(i)$ (Wz)(i)D(i,i) m $= \sum_{i=1}^{n} D(i,i) \neq (i)^{2} - \sum_{i=1}^{n} (w \neq) (i) \neq (i)$

 $= z^T D z - z^T \omega z = z^T L z$ M Since L = D - W.

Properties of the graph Laplacian

Lemma 6. Let L = D - W be the graph Laplacian corresponding to a symmetric matrix W with nonnegative entries. The following properties hold.

- (i) L is symmetric.
- (ii) L is positive semi-definite (i.e., $z^T L z \ge 0$ for all $z \in \mathbb{R}^m$).
- (iii) All eigenvalues of L are nonnegative, and the constant vector z = 1 is an eigenvector of L with eigenvalue $\lambda = 0$.

$$L1 = D1 - W1$$

(W1)(i) = $\sum_{j=1}^{m} W(i_j) 1 = D(i_j)$
=> $W1 = D1$.

Fiedler vector

Let v_1, v_2, \ldots, v_m be the eigenvectors of the graph Laplacian, with corresponding eigenvalues

$$0 = \lambda_1 \le \lambda_2 \le \dots \le \lambda_m.$$

Definition 7. The second eigenvector v_2 of the graph Laplacian L is called the *Fiedler vector*.

Theorem 8. The Fiedler vector $z = v_2$ solves the binary spectral clustering problem

Minimize
$$E(z)$$
 over $z \in \mathbb{R}^m$, subject to $\mathbf{1}^T z = 0$ and $||z||^2 = 1$.

Proof: Let 2 be a minimizer

Write
$$Z = a_1 V_1 + a_2 V_2 + \dots + a_m V_m$$

 $\left(V_1 = \frac{1}{5m}, \|V_1\| = 1\right)$
 $D = \frac{1}{7} = a_1 \frac{1}{V_1} + a_2 \frac{1}{2} \frac{1}{V_2} + \dots + a_m \frac{1}{5m}$
 $= a_1 \frac{m}{5m} = a_1 = 0$
 $I = \frac{1}{7} \frac{1}{7} = \sum_{i=1}^{m} a_i^2 = a_2^2 + a_3^2 + \dots + a_m^2$
 $E(Z) = \frac{2}{7} LZ = \frac{2}{7} L \sum_{i=2}^{m} a_i V_i$
 $= \frac{2}{7} \sum_{i=2}^{m} a_i L V_i$

$$= z^{T} \sum_{i=2}^{m} a_{i} \lambda_{i} \vee i$$

$$= \sum_{i=2}^{m} a_{i} \lambda_{i} z^{T} \vee i$$

$$= a_{i}$$

$$= \sum_{i=2}^{m} \lambda_{i} a_{i}^{2}$$

 $E(z) = \sum_{i=2}^{\infty} \lambda_i a_i^2 \ge \lambda_2 \sum_{i=2}^{\infty} a_i^2 = \lambda_2$ $\lambda_2 \le \lambda_3 \le \lambda_4 \le \dots \le \lambda_n$ Settim $z = V_2 \quad (a_2 = 1, a_3 = 0, a_4 = 0, \dots)$ yields $E(z) = \lambda_2$ XIIX

Example





Spectral embeddings

Let v_1, v_2, v_3, \ldots be the normalized eigenvectors of L, in order of increasing eigenvalues $0 = \lambda_1 \leq \lambda_2 \leq \cdots$. The spectral embedding corresponding to L is the map $\Phi: I_m \to \mathbb{R}^k$ (recall $I_m = \{1, 2, \ldots, m\}$ are the indices of our datapoints) given by

(9)
$$\Phi(i) = (v_1(i), v_2(i), \dots, v_k(i)).$$

Since the first eigenvector v_1 is the trivial constant eigenvector, it is also common to omit this to obtain the embedding

$$\Phi(i) = (v_2(i), v_3(i), \dots, v_{k+1}(i)).$$

There are other normalizations of the graph Laplacian that are commonly used, such as the symmetric normalization $L = D^{-1/2}(D-W)D^{-1/2}$, and the spectral embedding for a normalized Laplacian is defined analogously.

Spectral embeddings

The intuition behind the spectral embedding is encapsulated in the following simple result.

Proposition 9. If $A \subset I_m$ is a disconected component of the graph, which means that W(i, j) = 0 for all $i \in A$ and $j \in I_m \setminus A$, then the indicator function of A, denoted u_A , satisfies

 $Lu_A = (D - W)u_A = 0.$

Spectral embedding of MNIST



Figure 2: Example of spectral embeddings in the plane k = 2 of the 0, 1, and 2 digits of the MNIST dataset using the unnormalized L = D - W and symmetric normalized $L = D^{-1/2}(D - W)D^{-1/2}$ graph Laplacians.

Spectral clustering with more than 2 classes

Suppose we want to cluster the graph into k clusters. Spectral clustering proceeds as follows:

- 1. Perform a spectral embedding of the graph into \mathbb{R}^k (or sometimes \mathbb{R}^d where $d\approx k).$
- 2. Run your favorite clustering algorithm in the embedding space \mathbb{R}^d , such as k-means.

k-nearest neighbor graph

The Gaussian weights

$$W(i,j) = \exp\left(-\frac{\|x_i - x_j\|^2}{2\sigma^2}\right),\,$$

are not always useful in practice, since the matrix W is dense (all entries are nonzero), and the connectivity length σ is the same across the whole graph.

It is more common to use a k-nearest neighbor graph. Let $d_{k,i}$ denote the Euclidean distance between x_i and its k^{th} nearest Euclidean neighboring point from x_1, \ldots, x_m . A k-nearest neighbor graph uses the weights

$$W(i,j) = \begin{cases} 1, & \text{if } \|x_i - x_j\| \le \max\{d_{k,i}, d_{k,j}\}\\ 0, & \text{otherwise.} \end{cases}$$

The weights need not be binary, and can depend on $||x_i - x_j||$, similar to the Gaussian weights. The k-nearest neighbor graph weight matrix W is very sparse (most entries are zero), so it can be stored and computed with efficiently.

Exercises

The following proofs were omitted and are left to review as an exercise.

- 1. Proof of Theorem 2 (Convergence of k-means, Theorem 4.2 in lecture notes)
- 2. Proof of Lemma 6 (Properties of the graph Laplacian, Lemma 4.8 in lecture notes)
- 3. Proof of Proposition 9 (Spectral embeddings, Proposition 8.4 in lecture notes)

Spectral clustering in Python (.ipynb)